

CMSC423: Bioinformatic Algorithms, Databases and Tools Lecture 9

inexact alignment
dynamic programming, gapped
alignment

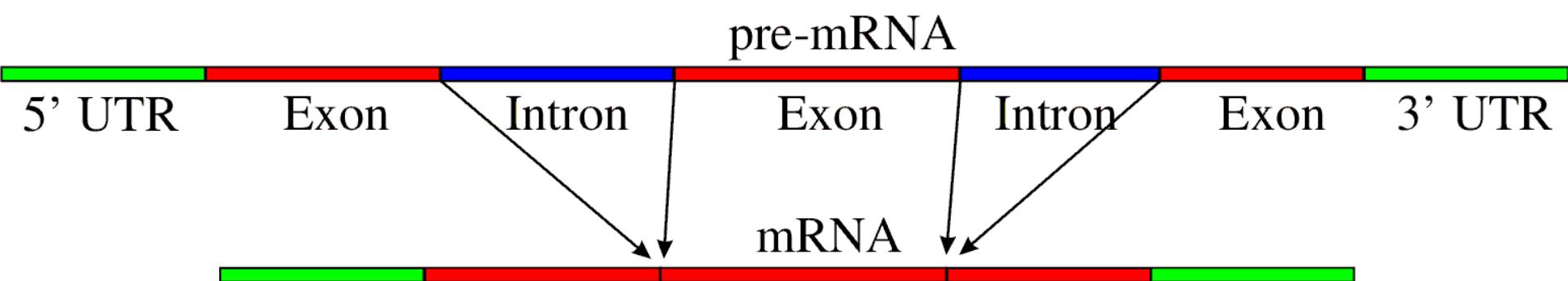
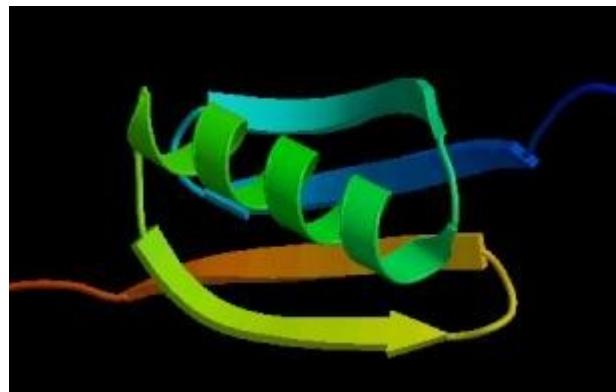
Inexact alignment

Inexact matching: why?

- Redundancy in genetic code: nucleotide sequence may differ, but proteins the same

S	Y	P	T	D
TCTTATCCTACTGAT				
TCATAACCCCACAGAC				

- Different amino-acid sequences still fold the same way: function unchanged (generally changing an amino-acid with a similar one doesn't affect protein function)
- Aligning ESTs (RNA sequences) to DNA need to account for gaps corresponding to exons
- Need to account for sequencing errors
- Read chap 6.1!!! (define: ortholog, paralog, homolog)



C	G	C	C	G	T	C	C	C	A	C	T	C	T	C	C	G	-	C	C	C	-	T	T	C	A	C	G	C	T	G
C	G	C	C	G	T	C	C	C	A	C	T	C	T	C	C	G	-	C	C	C	-	T	T	C	A	C	G	C	T	G
C	G	C	C	G	T	C	C	C	A	C	T	C	T	C	C	G	-	C	C	C	-	T	T	C	A	C	G	T	T	G
C	G	C	C	G	T	C	C	C	A	C	T	T	T	C	C	G	G	C	C	C	-	T	T	C	A	C	G	T	T	G
C	G	C	C	G	T	C	C	C	A	C	T	T	T	C	C	G	-	C	C	C	T	C	A	C	G	T	T	G		

Several hemoglobins

HBB_HUMAN	FFESFGDLSTPDAVMGNPKVKAHGKKVL-----GAFSDGLAHLDNLKGTF
HBB_HORSE	FFDSFGDLSNPGAVMGNPKVKAHGKKVL-----HSFGEGVHLDNLKGTF
HBA_HUMAN	YFPHF-DLS-----HGSAQVKGHGKKVA-----DALTNAVAHVDDMPNAL
HBA_HORSE	YFPHF-DLS-----HGSAQVKAHGKKVG-----DALTLAVGHLDLPGAL
MYG_PHYCA	KFDRFKHLKTEAEMKASEDLKKHGVTVL-----TALGAILKKKGHEAEL
GLB5_PETMA	FFPKFKGLTTADQLKKSADVRWHAERII-----NAVNDAVASMDDTEKMS
LGB2_LUPLU	LFSFLKGTSSEVP--QNNPELQAHAGKVFKLVYEAAIQLQVTGVVVTDATL * : : * . : : : :

From http://bioinfo.cnio.es/docus/courses/SEK2003Filogenias/seq_analysis/multiple.html

Warm-up – Longest Common Subsequence

- Given two strings of letters, identify longest string of letters that occurs, in the same order, in both strings

AG C GTAG
G C G A
GTCAG A

	A	G	C	G	T	A	G
G		1		1			1
T					1		
C			1				
A	1					1	
G		1		1			1
A	1					1	

- Find the longest chain of 1s, moving to the right and down

Dynamic programming

- Idea: re-use previously computed information
- $\text{LCS}[i,j]$ – longest common subsequence of strings $S1[1..i]$, $S2[1..j]$

					i		
	A	G	C	G	T	A	G
G		1		1			1
T					1		
C			1				
A	1					1	
G		1		1			1
A	1					1	

$\text{LCS}[i,j]$ is the maximum of:

1. if $S1[i] = S2[j]$
 $\text{LCS}[i-1, j-1] + 1$
- else
 $\text{LCS}[i - 1, j-1]$
- $\text{LCS}[i, j - 1]$

Computing the LCS table

Row 0 and column 0 easy to fill

Fill the rest column by column

Find the actual sequence:
trace-back pointers

	A	G	C	G	T	A	G
G	0	1	0	1	0	0	0
T	0	1					
C	0	1					
A	1	1					
G	0	2					
A	1	2					

	A	G	C	G	T	A	G
G	0	1	0	1	0	0	0
T	0	1	1	1	2	2	2
C	0	1	2	2	2	2	2
A	1	1	2	2	2	3	3
G	0	2	2	3	3	3	4
A	1	2	2	3	3	4	4

Extending to sequence alignment

AG-C-GTAG

-GTCAG-A-

- In LCS, mis-alignments were free
- What happens if we pay for our "mistakes"? (this also allows us to account for "similar" amino-acids)
 - Value[A, A] = 10
 - Value[A, G] = -5
 - Value[A, -] = -2
 - etc.
- The same dynamic programming algorithm works!

The recurrences

AG-C-GTAG
-GTCAG-A-

Score[i,j] is the maximum of:

1. Score[i-1, j-1] + Value[S1[i], S2[j]]

AG-C-G	AG-C-G
-GTCAG	-GTCAT

2. Score[i - 1, j] + Value[S1[i], -] (S1[i] aligned to gap)

AG-C-GT
-GTCAG-

3. Score[i, j - 1] + Value[-, S2[j]] (S2[j] aligned to gap)

AG-C-
-GTCA

The dynamic programming table

Score[i,j] is the maximum of:

1. Score[i-1, j-1] + Value[S1[i-1], S2[j-1]] (S1[i-1], S2[j-1] aligned)
2. Score[i – 1, j] + Value[S1[i], -] (S1[i] aligned to gap)
3. Score[i, j – 1] + Value[-, S2[j]] (S2[j] aligned to gap)

	-	A	G	C	G	T	A	G
-	0	-2	-4	-6	-8	-10	-12	-14
G	-2	-4	8	6				
T	-4	-6	6	4				
C	-6	-8	4	16				
A	-8							
G	-10							
A	-14							

$$\begin{aligned} \text{Value (A, A)} &= 10 \\ \text{Value (A, G)} &= -5 \\ \text{Value (A, -)} &= -2 \end{aligned}$$

Note: we only look at 3 adjacent boxes